



# CSC212

# Data Structure

## - Section FG

# Lecture 12

# Stacks and Queues

Instructor: Feng HU

Department of Computer Science

City College of New York

# Topics

- Stacks (Chapter 7)
- Queues (Chapter 8, Section 1 - 3)
- Priority Queues (Chapter 8, Section 4)
- References Return Values (Chapter 8, Section 5)

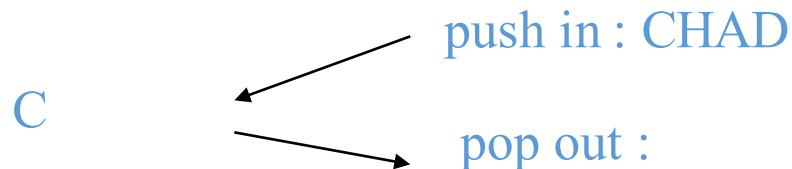
# Stacks and the STL **stack**

- Definition

- A stack is a data structure of *ordered* entries such that entries can be inserted and removed at only one end (call the top)

- LIFO

- A stack is a Last-In/First-Out data structure. Entries are taken out of the stack in the reverse order of their insertion



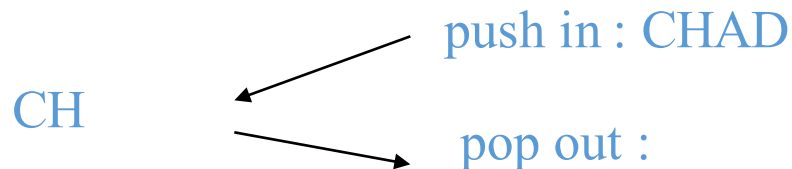
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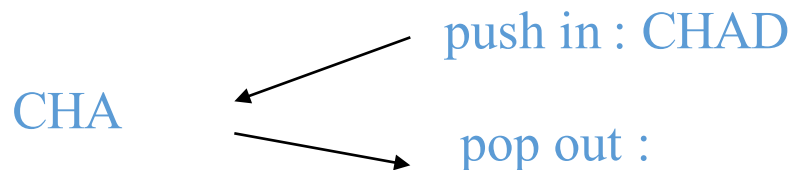
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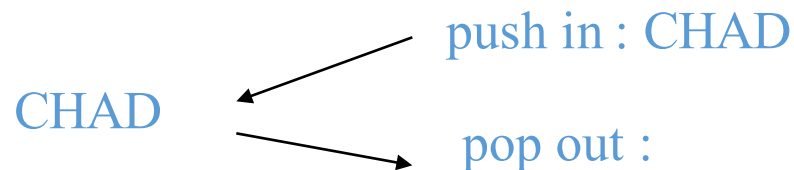
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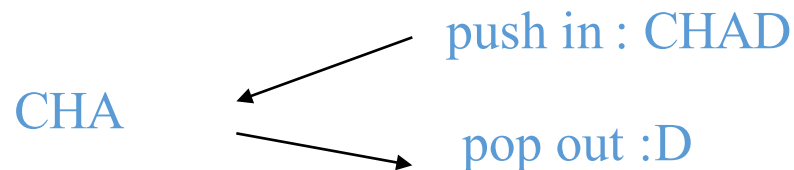
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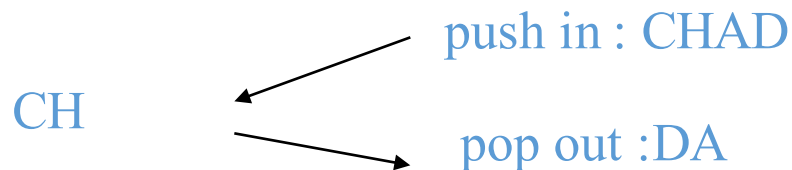
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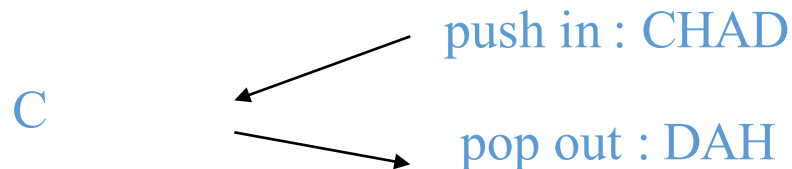
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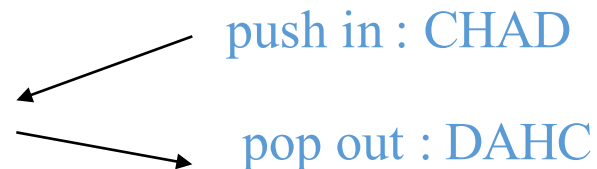
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# Stacks and the STL **stack**

- The STL stack class
  - a container class – holding many items
  - a template class – stack of integers, strings, ...
- How to use
  - `#include <stack>`
  - `stack<int> s1;`
- Implementation it ourselves! ([stack code](#))
  - fixed-size or dynamic array, or linked list?
  - STL typically uses dynamic array
  - Functions: push, pop, empty, size , **top**

# Queues and the STL queue

- Definition

- A queue is a data structure of *ordered* entries such that entries can only be inserted at one end (call the rear) and removed at the other end (call the front) – and the entry at the front of the queue is called the first entry

- FIFO

- A queue is a First-In/First-Out data structure. Entries are taken out of the queue in the same order that they were put into the queue

put in : CHAD      →      C      →      take out :

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put in : CHAD      →      HC      →      take out :

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put in : CHAD      →      DAHC      →      take out :

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put in : CHAD      →      DA      →      take out : CH

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put in : CHAD      → D      → take out : CHA

# Queues and the STL queue

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put in : CHAD      →                      →      take out : CHAD

# Queues and the STL `queue`

- The STL `queue` class
  - a container class – holding many items
  - a template class – queue of integers, strings, ...
- How to use
  - `#include <queue>`
  - `queue<char> q1;`
- Implementation it ourselves! ([queue code](#))
  - fixed-size or dynamic array, or linked list?
  - STL typically uses dynamic array
  - Functions: `push`, `pop`, `empty`, `size`, **`front`**

# Priority Queues

- A priority queue is a container class that allows entries to be retrieved according to some specified priority levels.
  - The highest priority entry is removed first
  - Entries with equal priority can be removed according some criterion e.g. FIFO as an queue.
- STL `priority_queue<Item>` template class
  - `#include <queue>`
  - `priority_queue<int> q2;`
  - Functions `push`, `pop`, `empty`, `size` , **top** (*not front!*)
  - *Several ways to specify priority (p. 411)*

# Reference Return Values for the stack, queue, and priority queue classes

- In STL, the top (for stack) and front (for queue) functions have reference return values, e.g. in stack class definition:

- Item& top ();
- const Item& top() const;

- Can be used to change the top item

- If we declare
  - stack<int> b;
  - const stack<int> c;
- Which ones are correct? =>

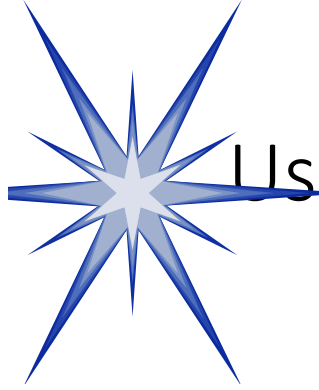
1. int i = b.top(); ✓

2. b.push(i); ✓

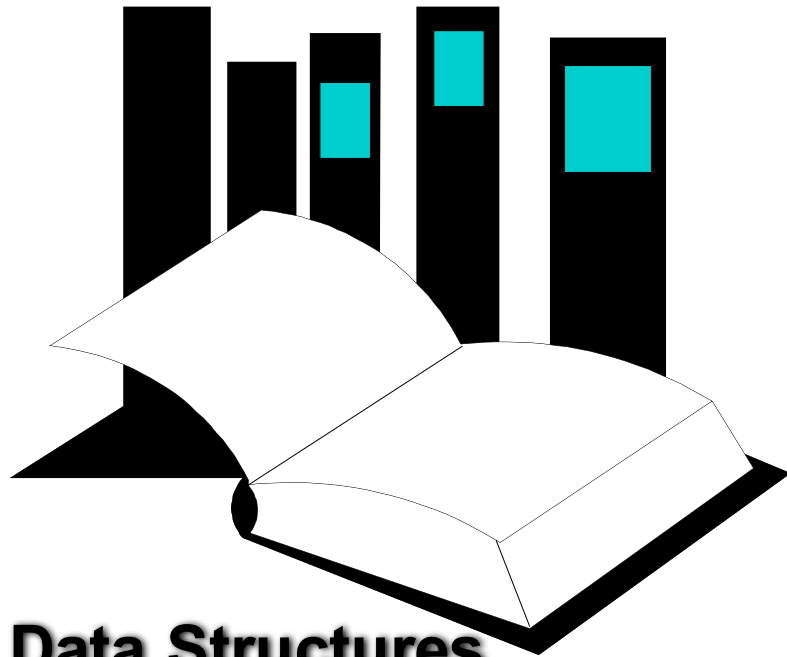
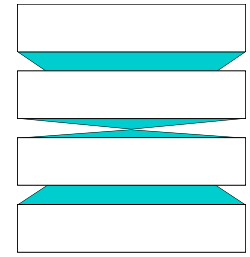
3. b.top() = 18; ✓

4. c.top() = 18; ✗

5. b.push(c.top()); ✓



# Using a Stack



**Data Structures  
and Other Objects  
Using C++**

- Chapter 7 introduces the stack data type.
- Several example applications of stacks are given in that chapter.
- This presentation shows another use called backtracking to solve the N-Queens problem.

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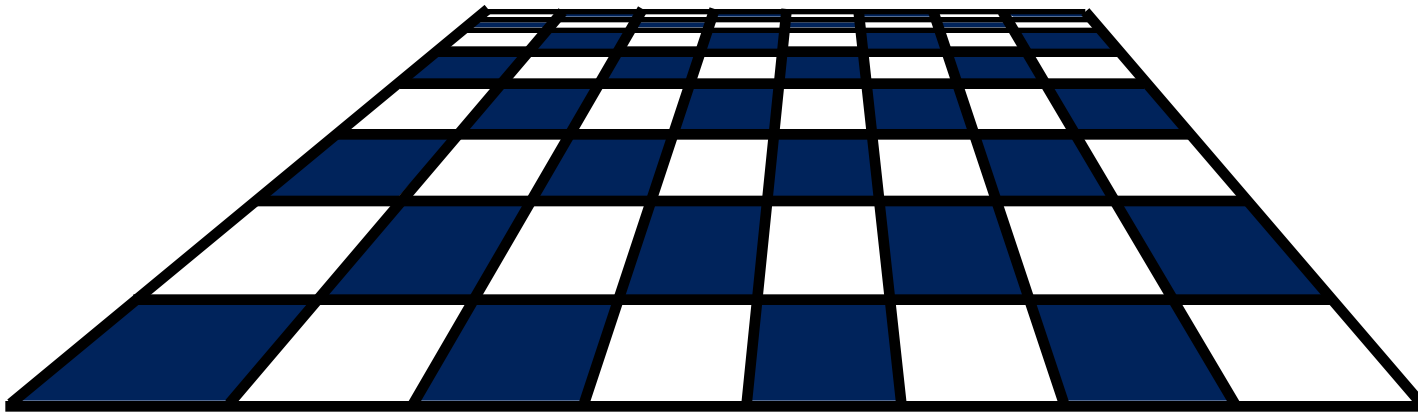
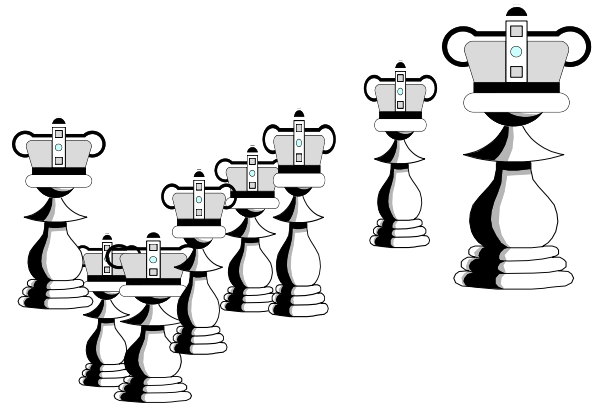
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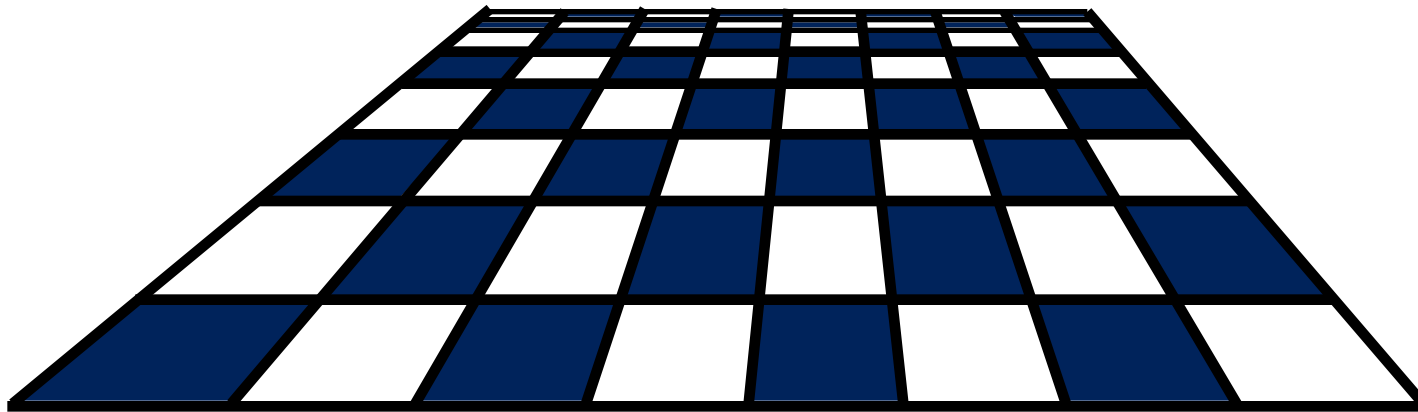
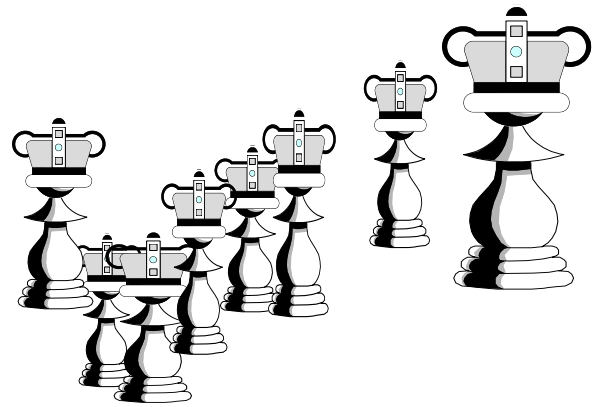
# The N-Queens Problem

- Suppose you have 8 chess queens...
- ...and a chess board



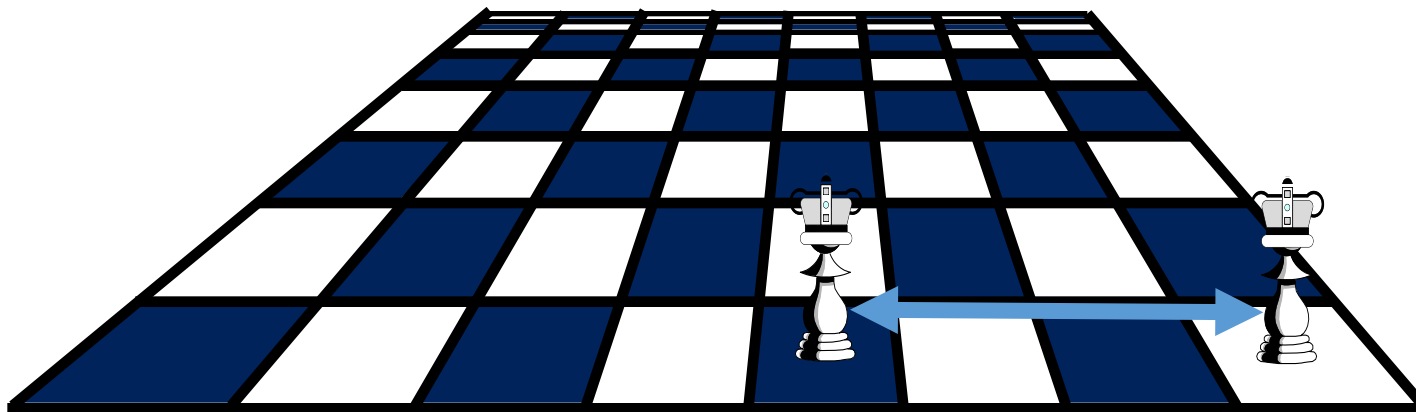
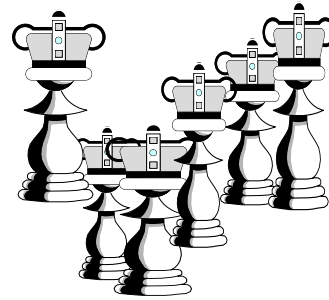
# The N-Queens Problem

*Can the queens be placed on the board so that no two queens are attacking each other*



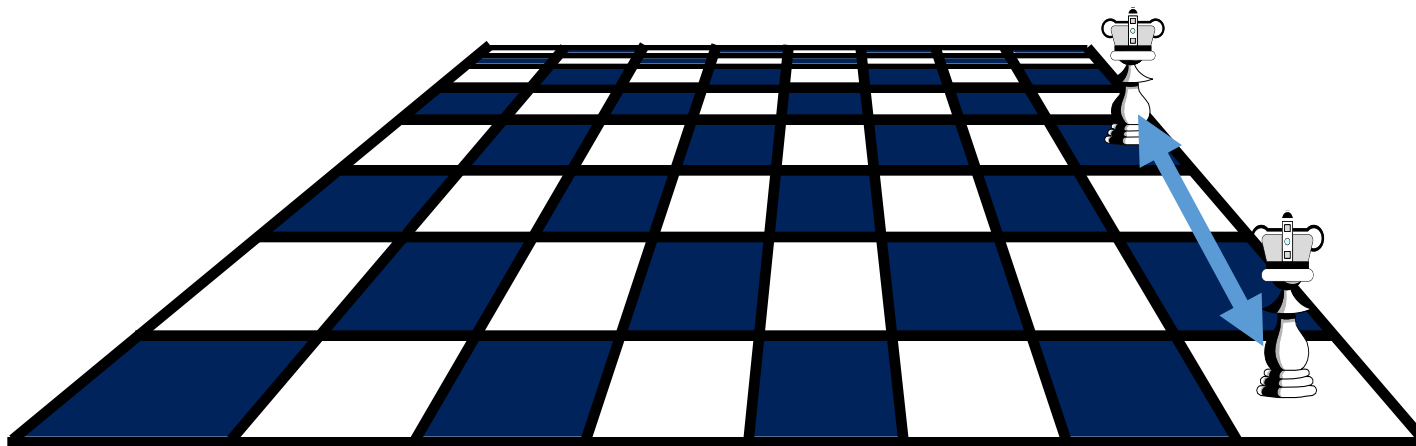
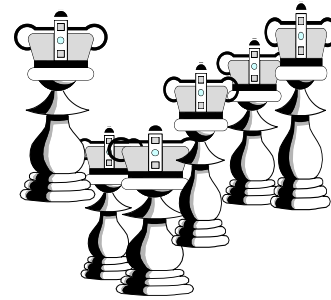
# The N-Queens Problem

Two queens are not allowed in the same row...



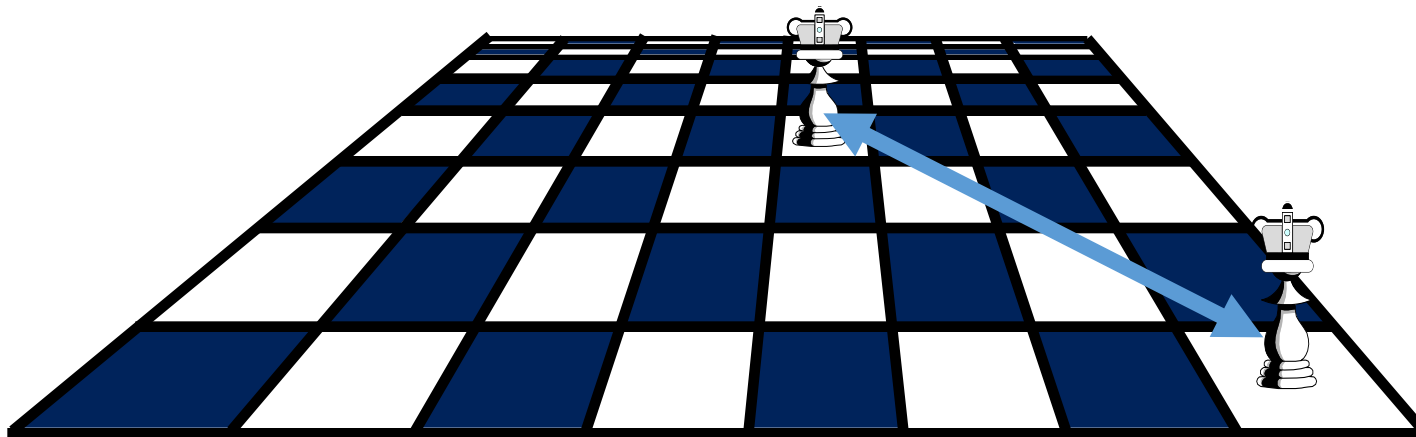
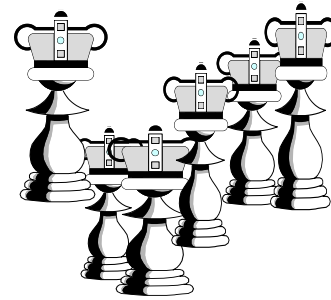
# The N-Queens Problem

Two queens are not allowed in the same row, or in the same column...



# The N-Queens Problem

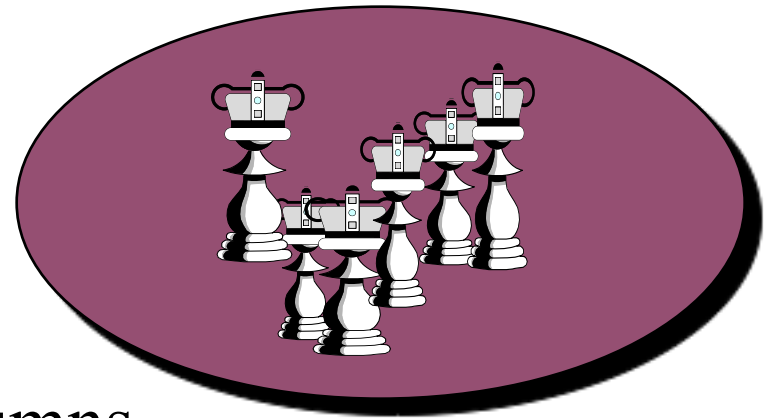
Two queens are not allowed in the same row, or in the same column, or along the same diagonal.



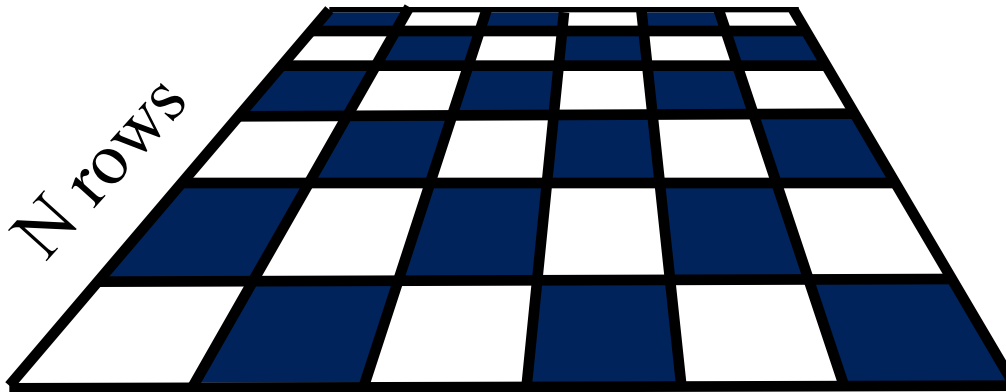
# The N-Queens Problem

The number of queens, and the size of the board can vary.

N Queens

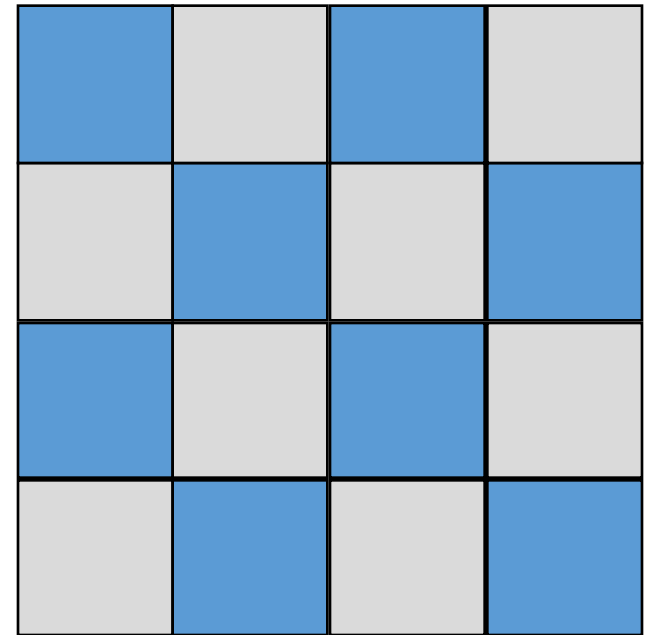
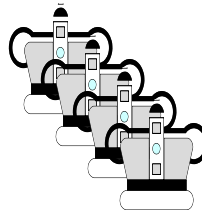
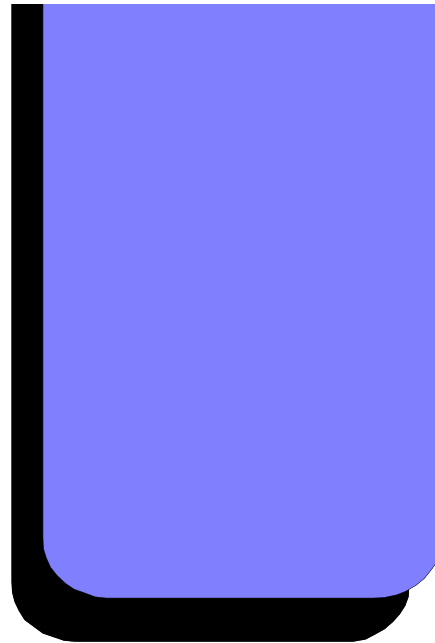


N columns



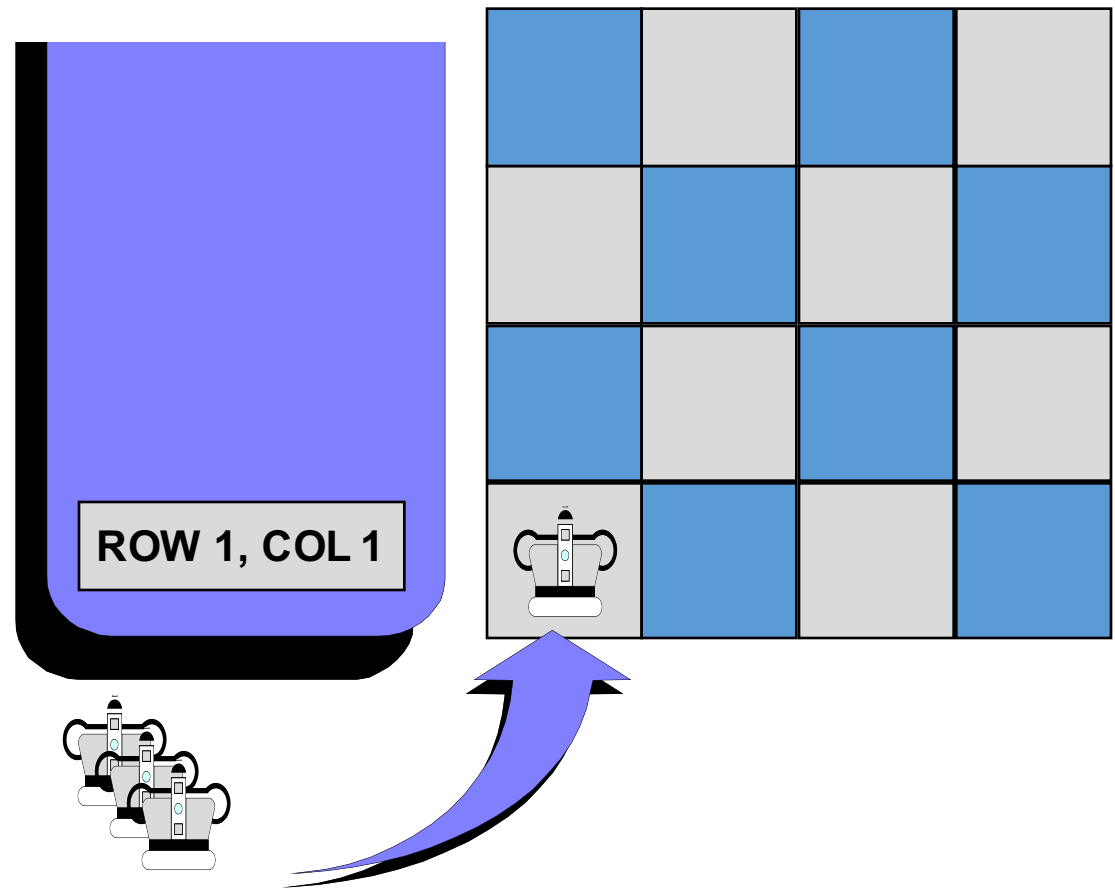
# How the program works

The program uses a stack to keep track of where each queen is placed.



# How the program works

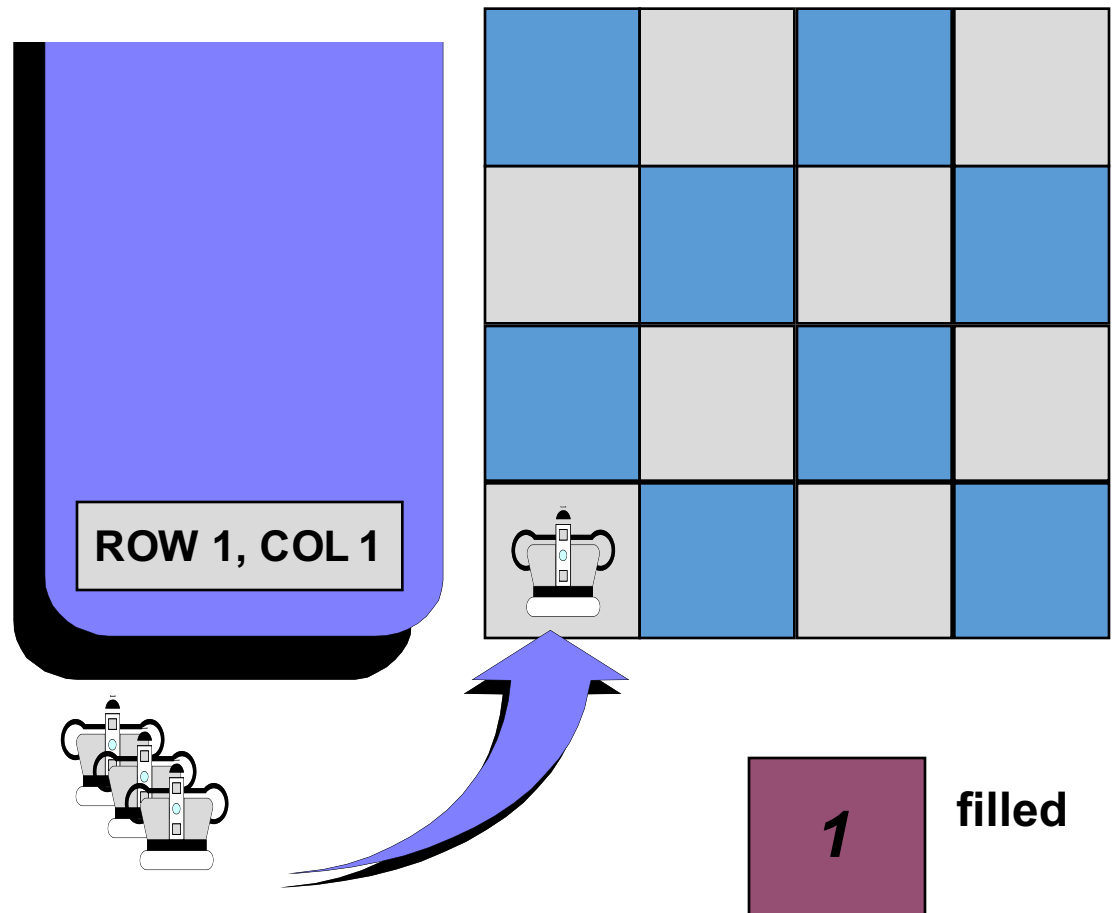
Each time the program decides to place a queen on the board, the position of the new queen is stored in a record which is placed in the stack.





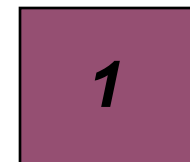
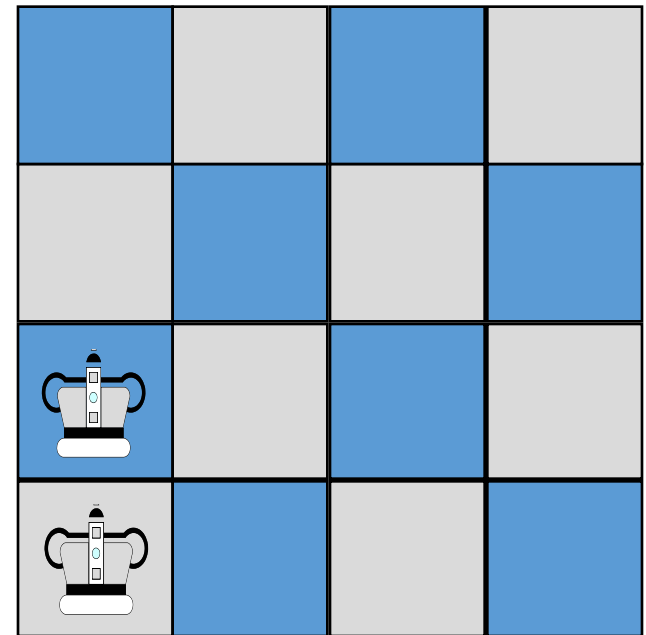
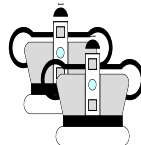
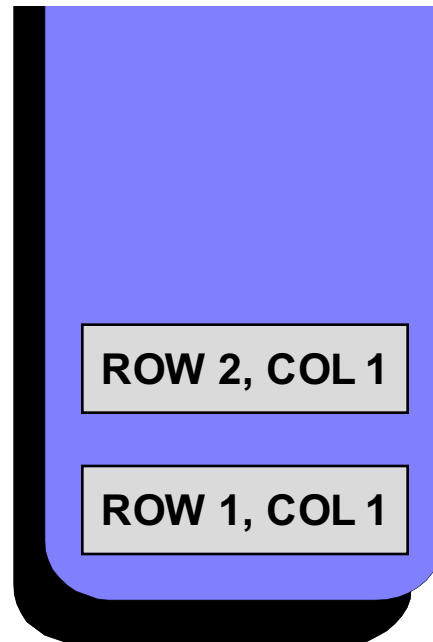
# How the program works

We also have an integer variable to keep track of how many rows have been filled so far.



# How the program works

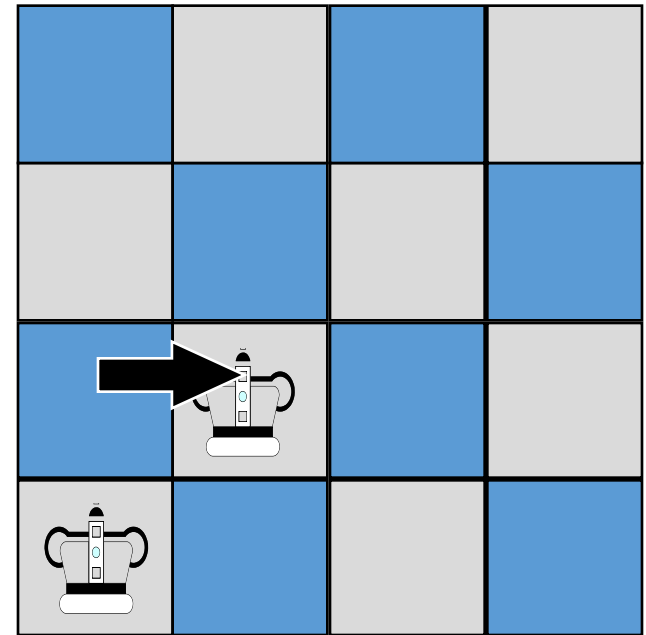
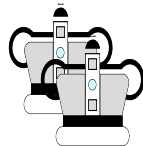
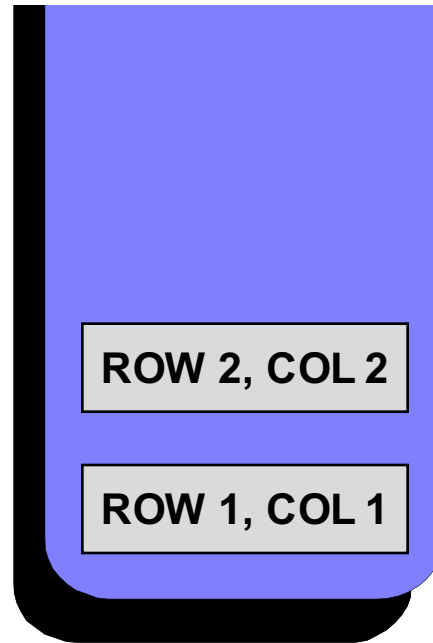
Each time we try to place a new queen in the next row, we start by placing the queen in the first column...



filled

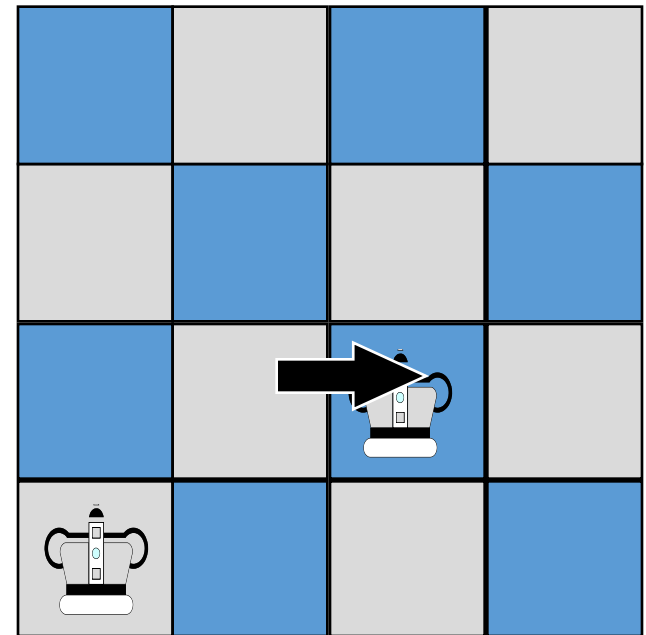
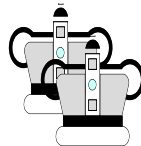
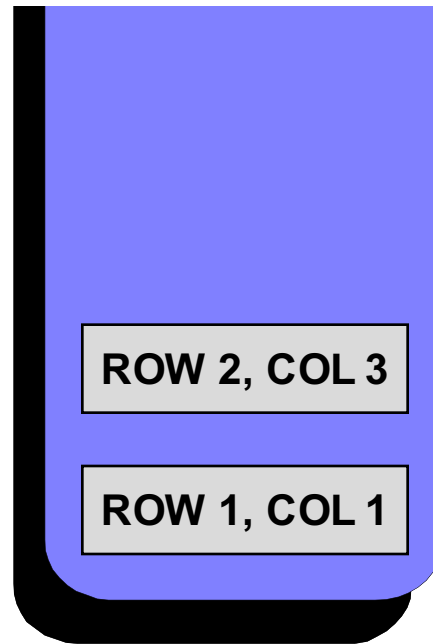
# How the program works

...if there is a conflict with another queen, then we shift the new queen to the next column.



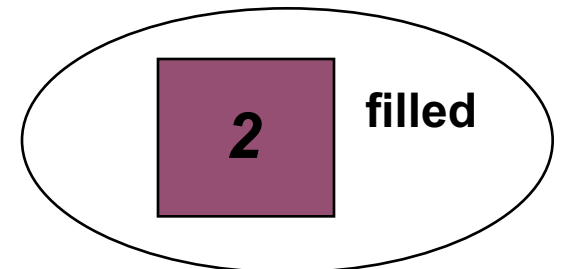
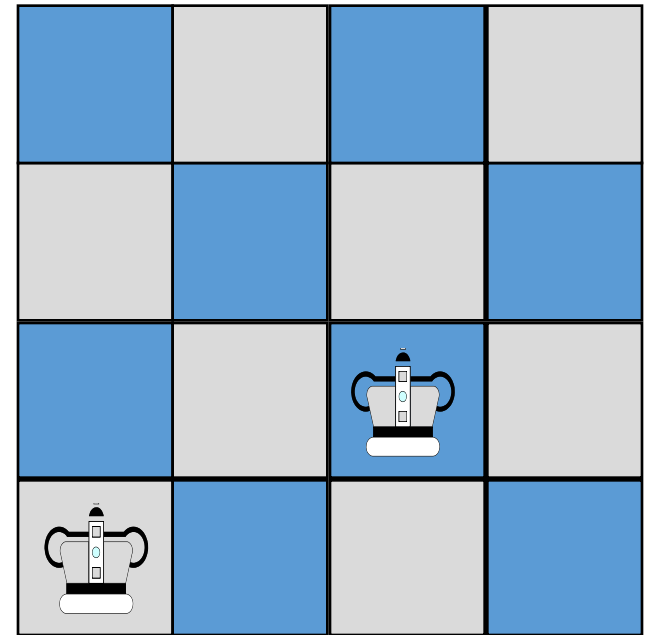
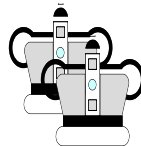
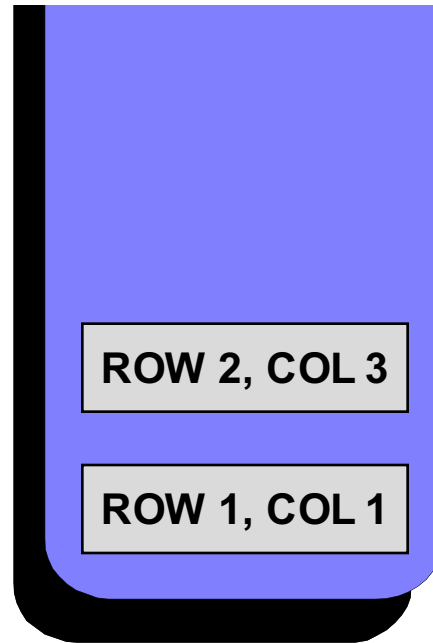
# How the program works

If another conflict occurs, the queen is shifted rightward again.



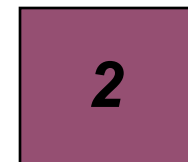
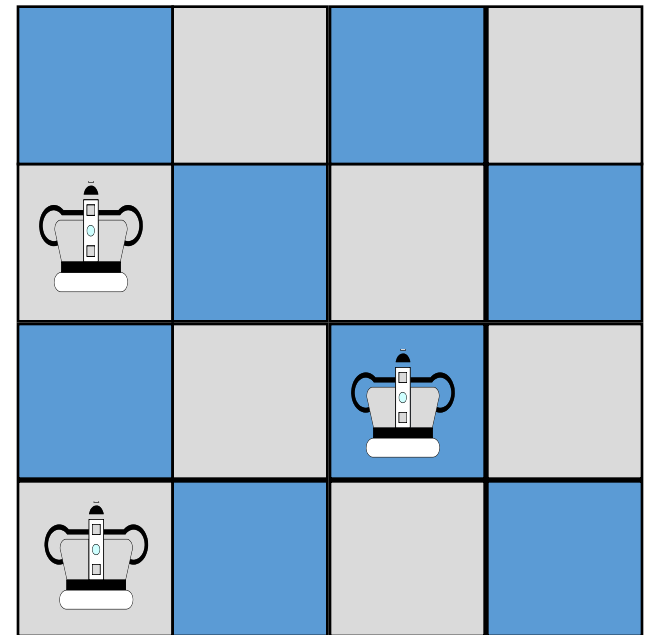
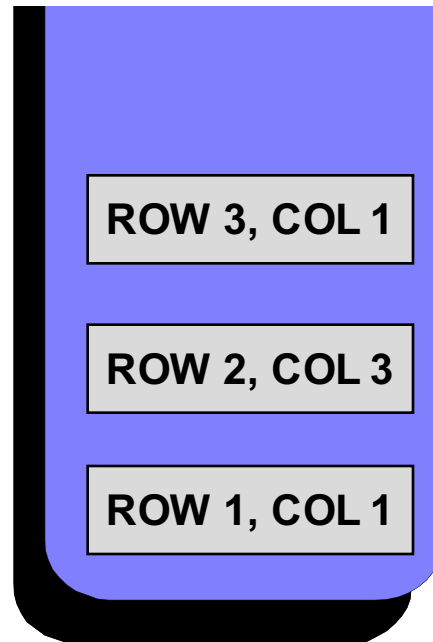
# How the program works

When there are no conflicts, we stop and add one to the value of filled.



# How the program works

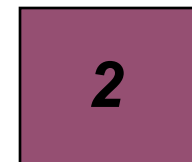
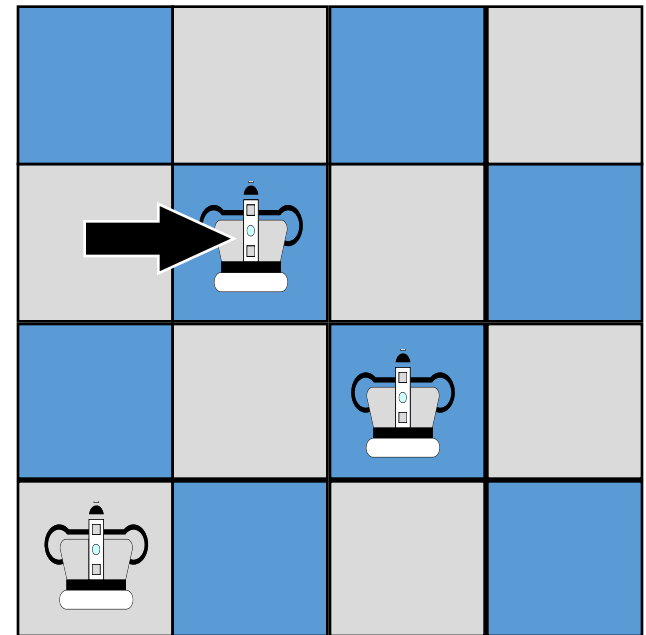
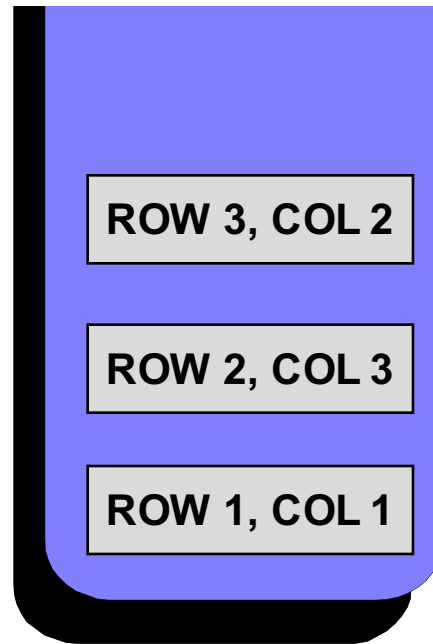
Let's look at the third row. The first position we try has a conflict...



filled

# How the program works

...so we shift to column 2. But another conflict arises...

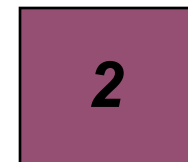
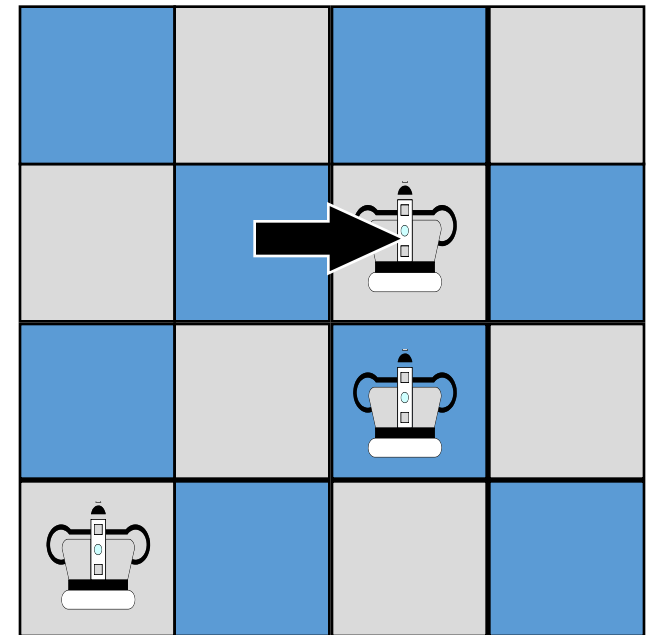
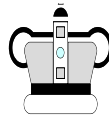
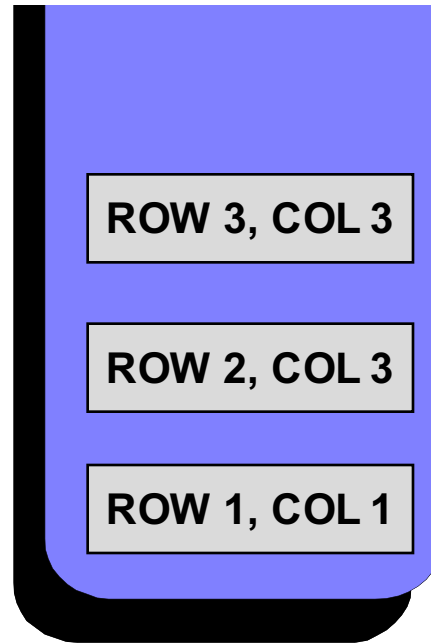


filled

# How the program works

...and we shift to the third column.

Yet another conflict arises...

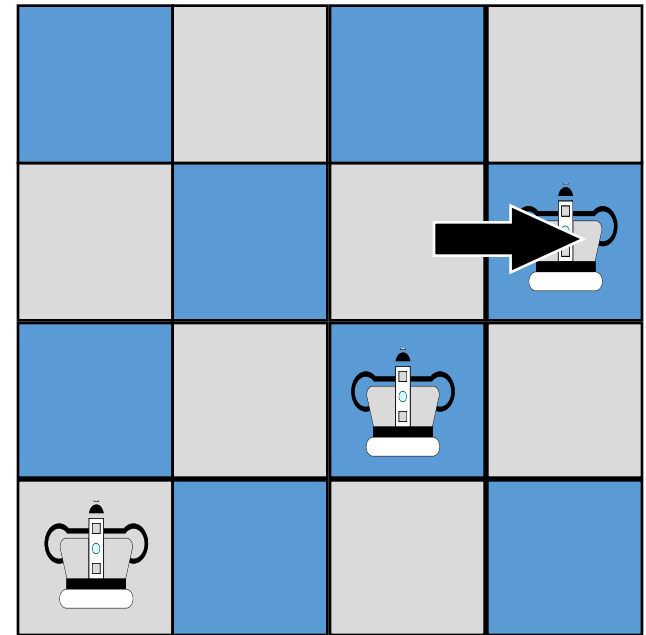
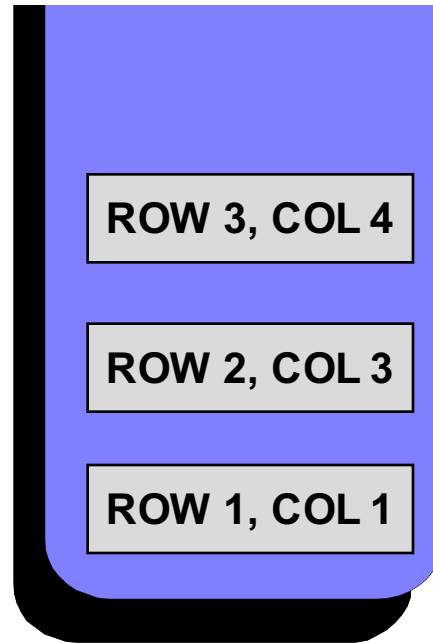


filled



# How the program works

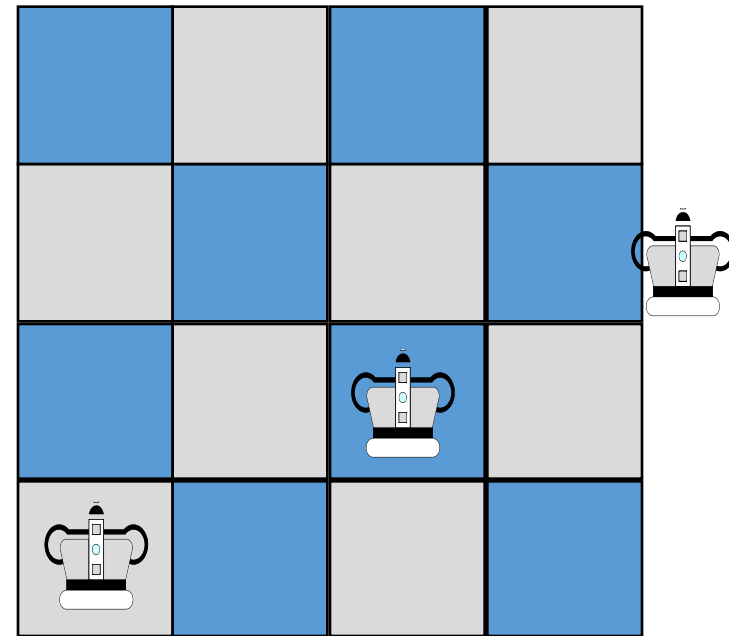
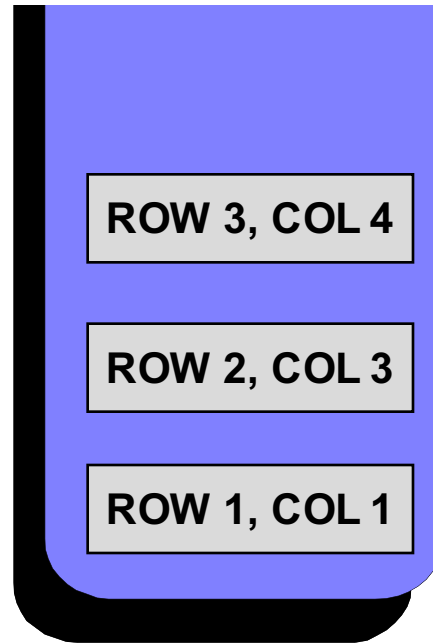
...and we shift to column 4. There's still a conflict in column 4, so we try to shift rightward again...



filled

# How the program works

...but there's nowhere else to go.

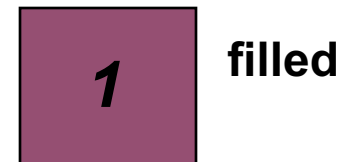
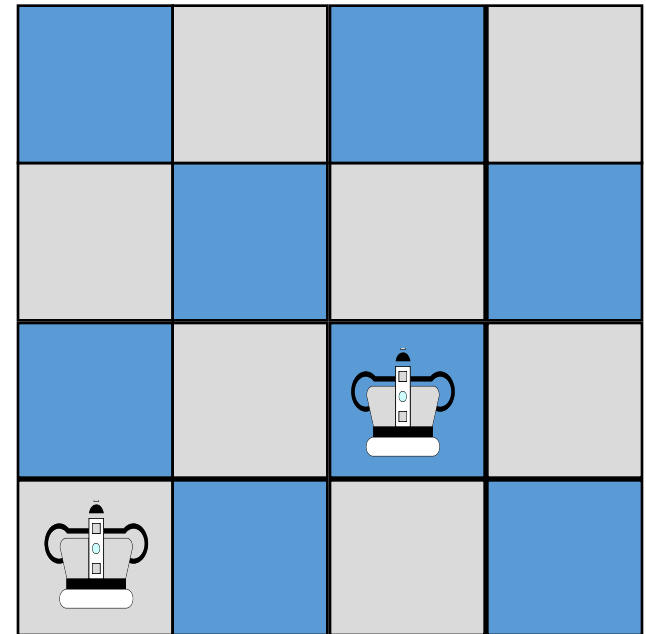
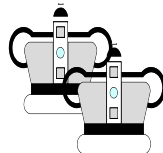
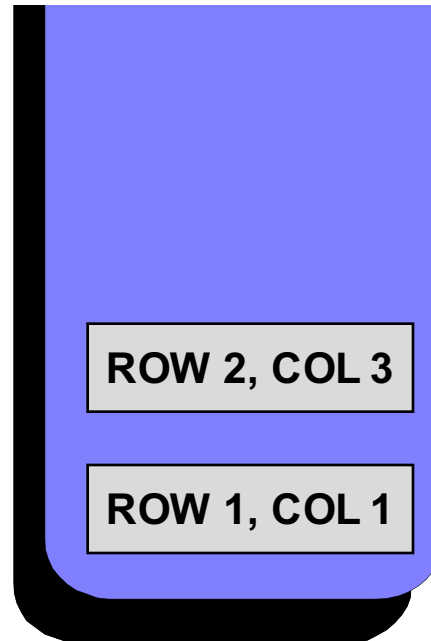


**2** filled

# How the program works

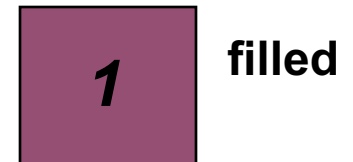
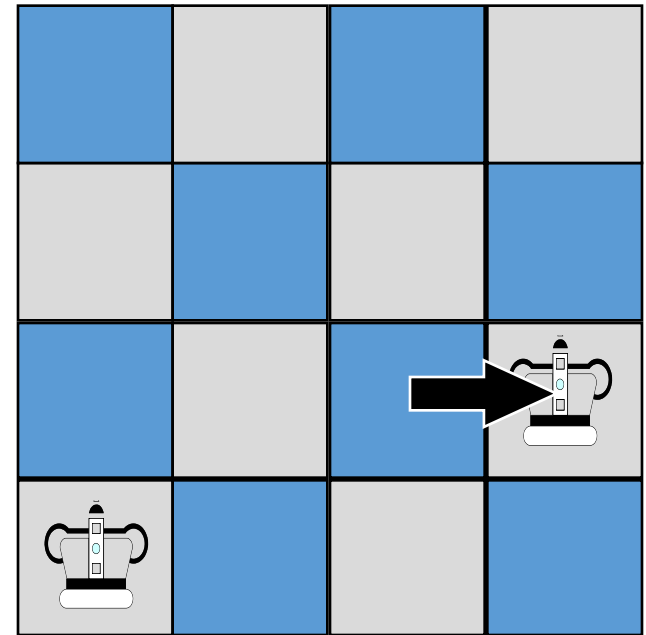
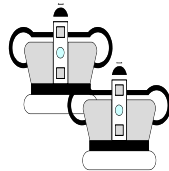
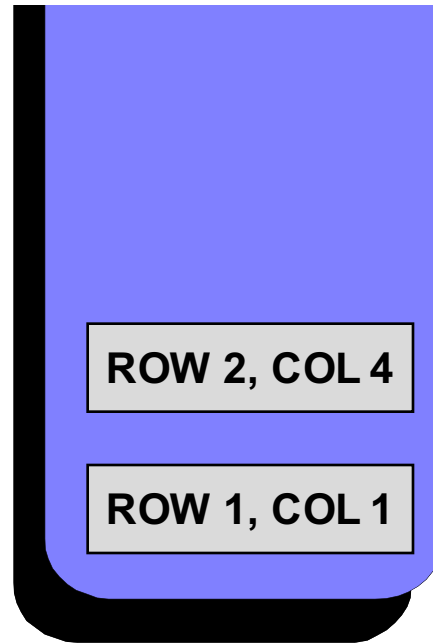
When we run out of room in a row:

- pop the stack,
- reduce filled by 1
- and continue working on the previous row.



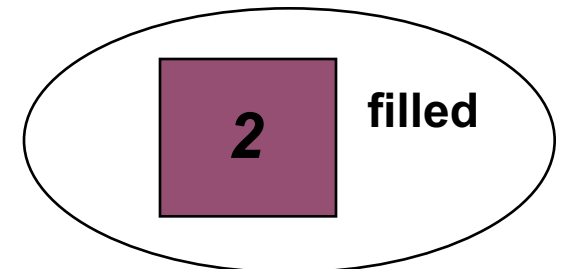
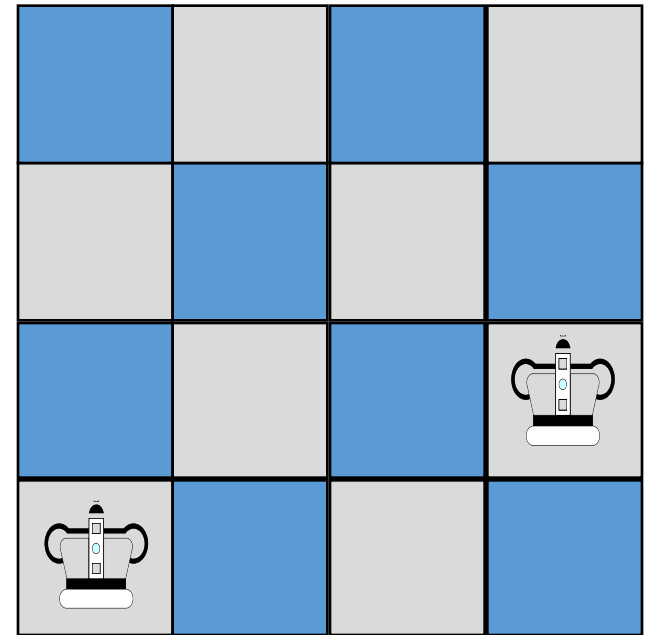
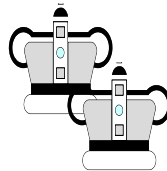
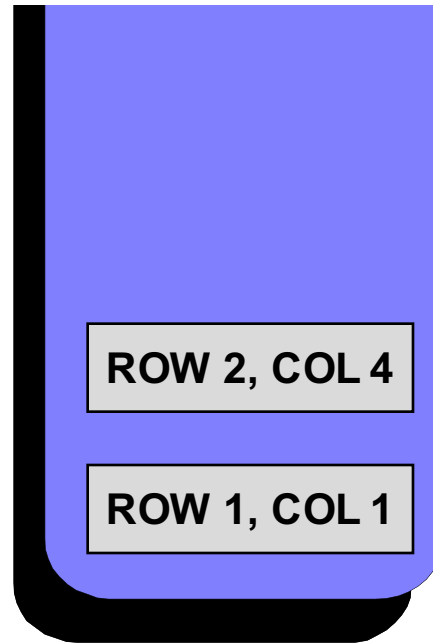
# How the program works

Now we continue working on row 2, shifting the queen to the right.



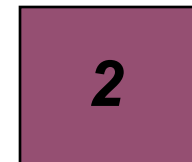
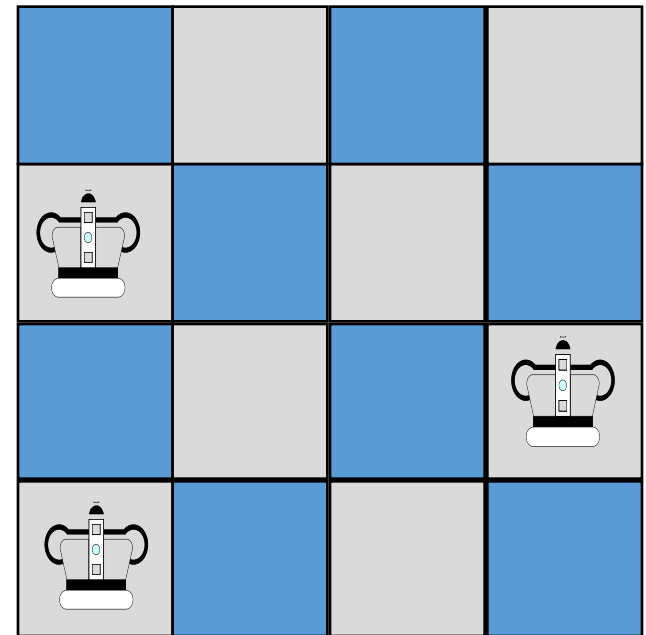
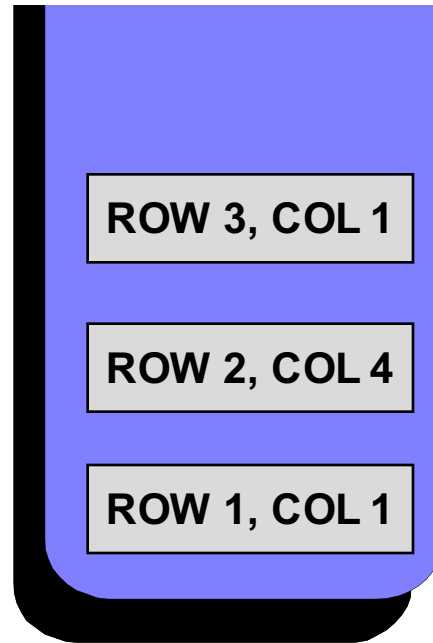
# How the program works

This position has no conflicts, so we can increase filled by 1, and move to row 3.



# How the program works

In row 3, we start again at the first column.



filled

# Pseudocode for N-Queens

- ① Initialize a stack where we can keep track of our decisions.
- ② Place the first queen, pushing its position onto the stack and setting filled to 0.
- ③ repeat these steps
  - if there are no conflicts with the queens...
  - else if there is a conflict and there is room to shift the current queen rightward...
  - else if there is a conflict and there is no room to shift the current queen rightward...

# Pseudocode for N-Queens

③ repeat these steps

- if there are no conflicts with the queens...

Increase filled by 1. If filled is now  $N$ , then the algorithm is done. Otherwise, move to the next row and place a queen in the first column.



# Pseudocode for N-Queens

## ③ repeat these steps

- if there are no conflicts with the queens...
- else if there is a conflict and there is room to shift the current queen rightward...

Move the current queen rightward,  
adjusting the record on top of the stack  
to indicate the new position.

# Pseudocode for N-Queens

## ③ repeat these steps

- if there are no conflicts with the queens...
- else if there is a conflict and there is room to shift the current queen rightward...
- else if there is a conflict and there is no room to shift the current queen rightward...

### Backtrack!

Keep popping the stack, and reducing filled by 1, until you reach a row where the queen can be shifted rightward. Shift this queen right.

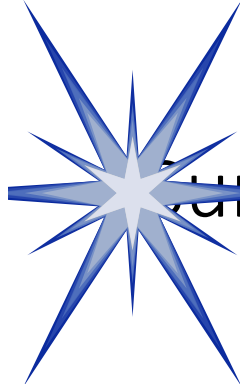
# Pseudocode for N-Queens

## ③ repeat these steps

- if there are no conflicts with the queens...
- else if there is a conflict and there is room to shift the current queen rightward...
- else if there is a conflict and there is no room to shift the current queen rightward...

### Backtrack!

Keep popping the stack, and reducing filled by 1, until you reach a row where the queen can be shifted rightward. Shift this queen right.



## Summary of stack for backtracking

- Stacks have many applications.
- The application which we have shown is called backtracking.
- The key to backtracking: Each choice is recorded in a stack.
- When you run out of choices for the current decision, you pop the stack, and continue trying different choices for the previous decision.

# Summary and Homework

- Stacks (Read Chapter 7)
  - Self-Test: 1-5, 13-18
- Queues (Read Sections 8.1 – 8.3)
  - Self-Test: 1-5, 10,18-21
- Priority Queues (Read Section 8.4)
  - Self-Test: 25-27
- References Return Values (Read Section 8.5 and p. 302 in Chapter 6)
  - Self-Test: class note